

# An Infrastructure for Implementing Compilers for Concurrent Abstract State Machine Languages

Kristian Magnani

Mariza A. S. Bigonha

Roberto S. Bigonha

Fabíola F. Oliveira

Vladimir O. Di Iorio

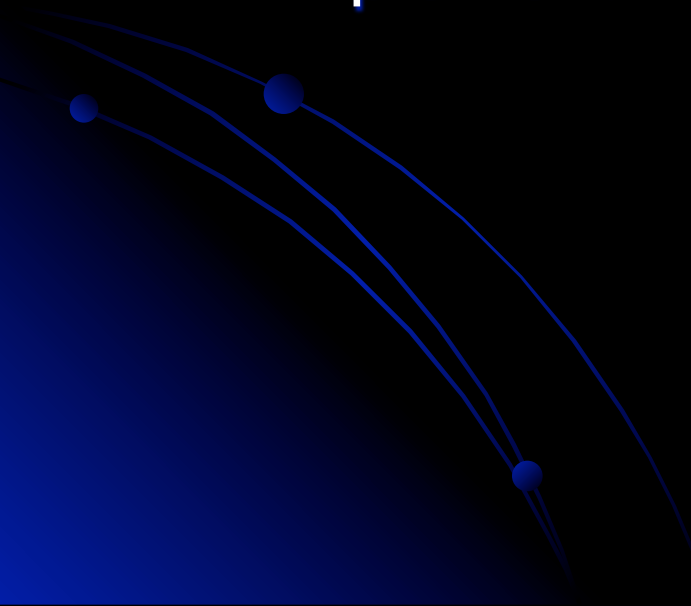
# Introduction

- Abstract State Machines (ASM)
  - formal semantic method created by Yuri Gurevich
  - provide operational semantics for algorithms
  - used in:
    - Semantics of programming languages
    - Distributed systems
    - Architectures (hardware and software)
    - etc

# Introduction

## This work:

- describes a new model for specifying concurrent systems based on ASM
- proposes an infrastructure, called MIR, to implement ASM-like languages

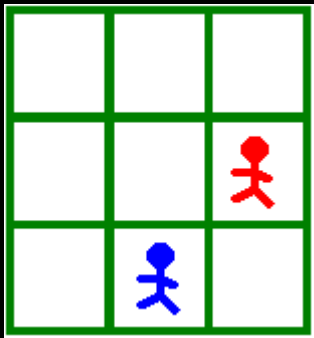


# Abstract State Machines (ASM)

- Machines whose states are algebraic structures, represented by *functions*
- Basic transition rules:
  - function update
  - conditional constructor
  - block constructor (parallel execution)
- Runs:
  - sequences of states
  - next state generated by the application of the transition rule over the previous state

# Abstract State Machines (ASM)

Example:



State:

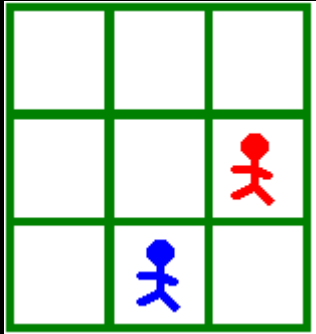
`lin (RED) = 2`

`col (RED) = 3`

`lin (BLUE) = 3`

`col (BLUE) = 2`

# Abstract State Machines (ASM)



`lin(RED) = 2`

`col(RED) = 3`

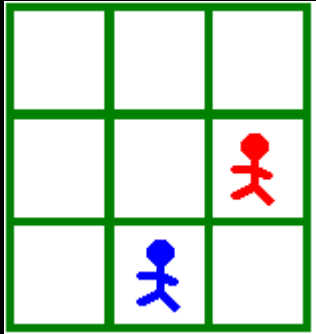
`lin(BLUE) = 3`

`col(BLUE) = 2`

Transition rule:

```
if col(RED) > 1 then
  col(RED) := col(RED) - 1
end
```

# Abstract State Machines (ASM)



`lin(RED) = 2`

`col(RED) = 3`

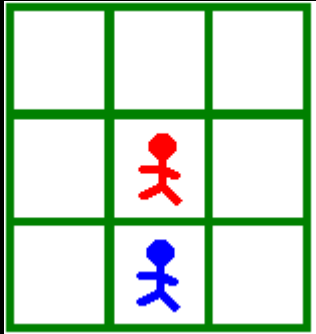
`lin(BLUE) = 3`

`col(BLUE) = 2`

```
if col(RED) > 1 then
  col(RED) := col(RED) - 1
end
```

Run: ...

# Abstract State Machines (ASM)



$\text{lin}(\text{RED}) = 2$

$\text{col}(\text{RED}) = 2$

$\text{lin}(\text{BLUE}) = 3$

$\text{col}(\text{BLUE}) = 2$

```
if col(RED) > 1 then
  col(RED) := col(RED) - 1
end
```

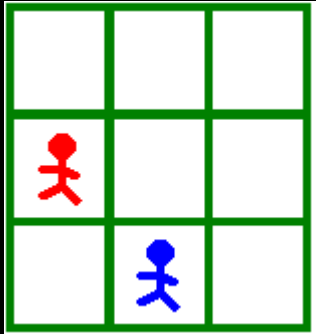
Run:

Step 1 ...





# Abstract State Machines (ASM)



`lin(RED) = 2`

`col(RED) = 1`

`lin(BLUE) = 3`

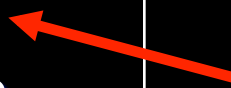
`col(BLUE) = 2`

```
if col(RED) > 1 then
  col(RED) := col(RED) - 1
end
```

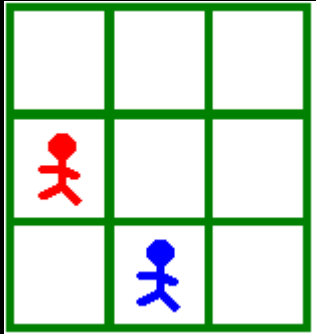
Run:

Step 1

Step 2 ...



# Abstract State Machines (ASM)



`lin(RED) = 2`

`col(RED) = 1`

`lin(BLUE) = 3`

`col(BLUE) = 2`

```
if col(RED) > 1 then
  col(RED) := col(RED) - 1
end
```

Run:

Step 1

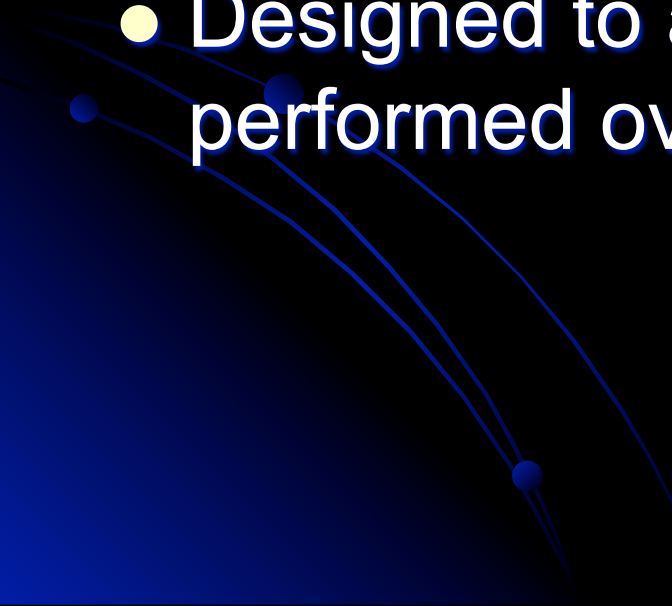
Step 2

Step 3 ...

# Abstract State Machines (ASM)

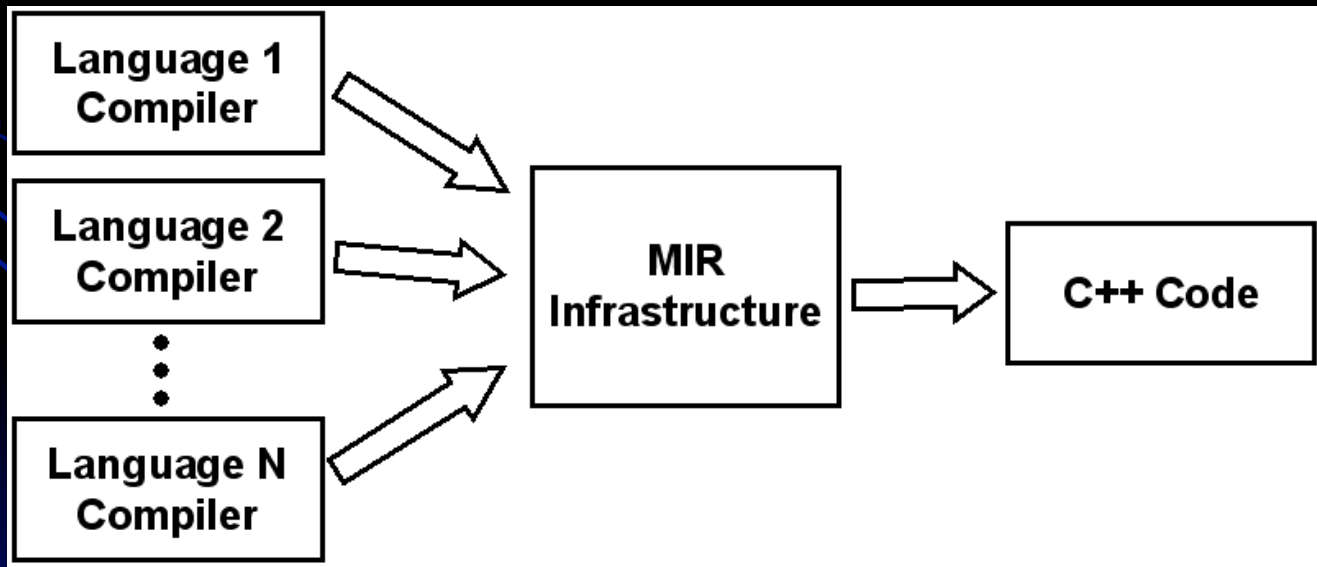
- Advantages:
    - Precision
    - Simplicity
    - Executability
    - Scalability
    - Generality
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# MIR

- Infrastructure for ASM implementation
  - Implements a concurrency model based on Lamport's concept of delayed knowledge
  - Designed to allow optimizations to be performed over ASM specifications
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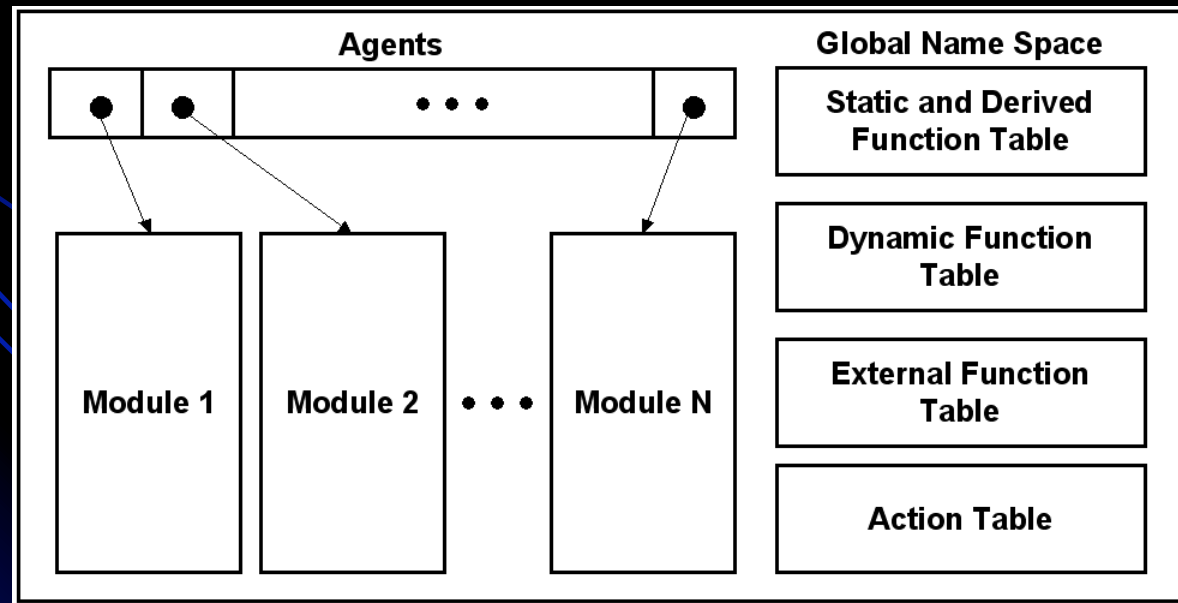
# The MIR Architecture

- Motivation:
  - general infrastructure used as the basis of compilers aiming at different ASM oriented languages;



# Agents, Modules and Other Elements

- A MIR specification:
  - a set of agents, each of them of a given type
  - a common Global Name Space, which is accessible by each agent.



# The MIR Approach for Concurrency

- An agent of MIR :

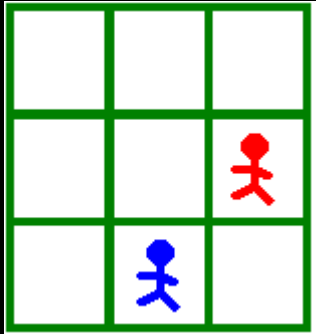
- **Autonomy**: there is a transition rule associated with each agent, giving it an autonomous and independent behaviour from other agents;
- **Situatedness**: every agent is immersed in a common global context where they can interact among each other, which are referenced by their names or by the special word self;
- **Proactivity**: an agent has the freedom to call actions over other agents.

# The MIR Approach for Concurrency

- Delayed knowledge (Lamport):
  - agents may store copies of the global state which may not always be kept up to date.
- Systems with agents of different speed of execution
- Unpredictable delay between the moment some changes are performed by one agent and the moment when these changes are perceived by another agent



# The MIR Approach for Concurrency



`lin(RED) = 2`

`col(RED) = 3`

`lin(BLUE) = 3`

`col(BLUE) = 2`

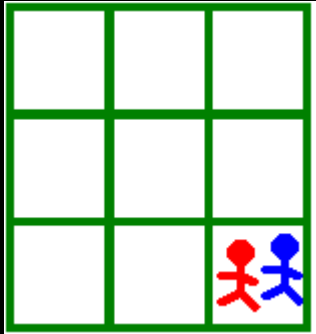
Rule for RED:

```
if isFree (lin(self)+1,  
          col(self)) then  
  lin(self) := lin(self)+1  
end
```

Rule for BLUE:

```
if isFree (lin(self),  
          col(self)+1) then  
  col(self) := col(self)+1  
End
```

# The MIR Approach for Concurrency



`lin(RED) = 3`

`col(RED) = 3`

`lin(BLUE) = 3`

`col(BLUE) = 3`

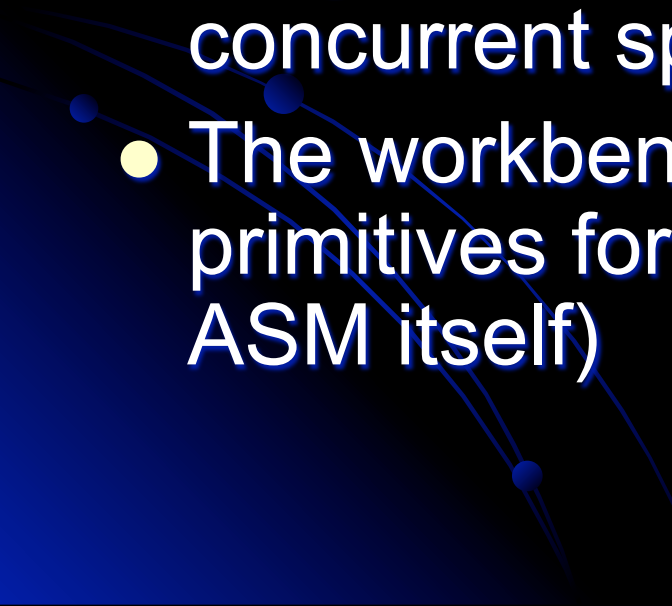
Rule for RED:

```
if isFree (lin(self)+1,  
           col(self)) then  
    lin(self) := lin(self)+1  
end
```

Rule for BLUE:

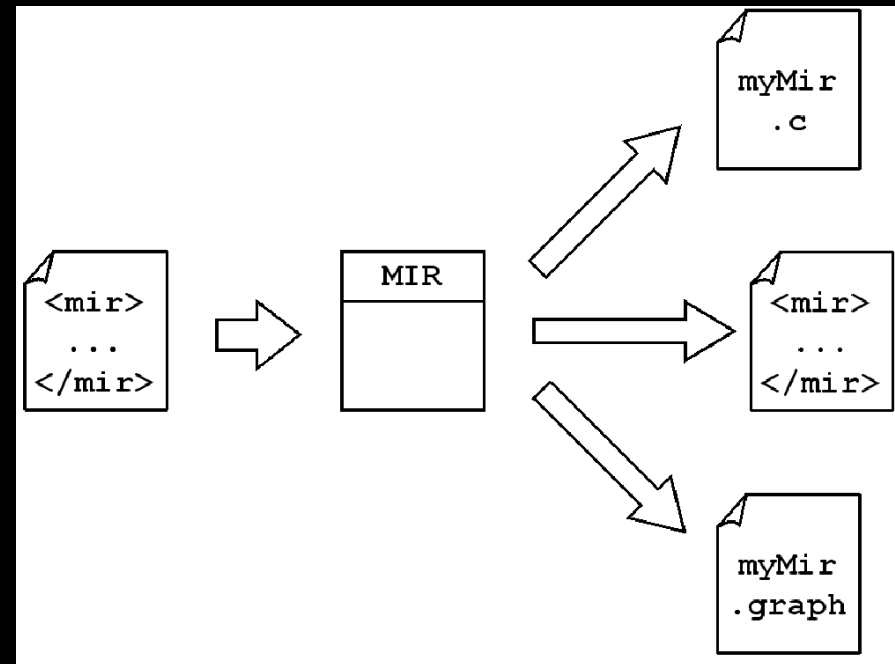
```
if isFree (lin(self),  
           col(self)+1) then  
    col(self) := col(self)+1  
End
```

# The MIR Approach for Concurrency

- The proposed model does not provide any primitive for synchronization between agents
  - All communication must be explicitly programmed by the designer of an ASM concurrent specification
  - The workbench may provide a library with primitives for synchronization (written in ASM itself)
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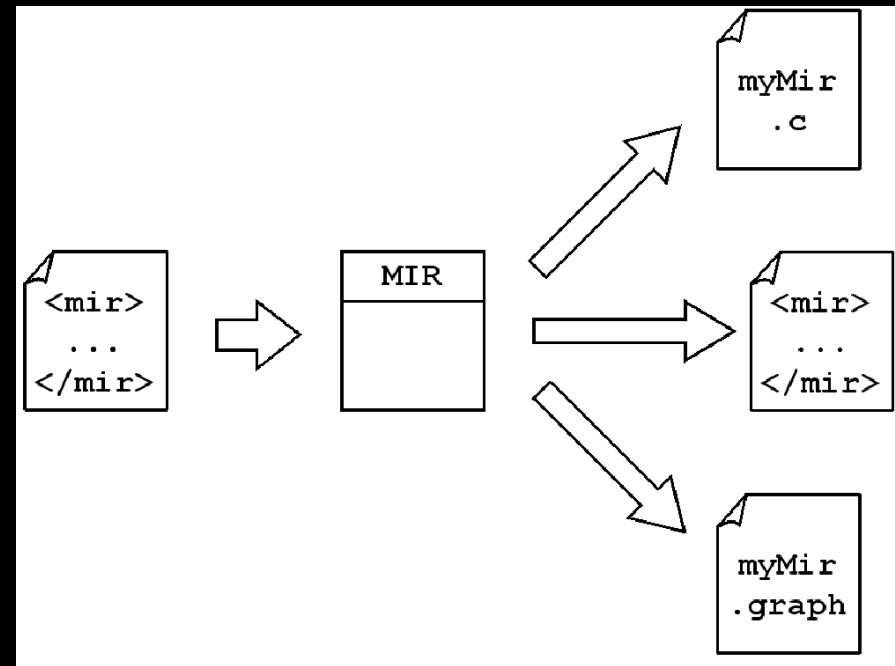
# Highlights of MIR Implementation

- **Serialization:** The MIR representation of an ASM specification can be saved in persistent store as a XML file
- **Compilation:** It can be compiled into C++ Code
- **Direct Execution:** MIR can be also interpreted



# Highlights of MIR Implementation

- **Visualization:** It is also possible to obtain a visual representation of a MIR specification through the generation of its description in the DOT language of the GraphViz software.



# MIR and Optimization

- There are some optimization opportunities that are particular to the ASM model, and therefore they are not performed by the existing C++ compilers
- In order to address this special situation, it is under development a framework that provides the proper environment to plugin specific MIR optimizations

# Conclusions

- Contributions of the proposed infrastructure:
  - It can be used to implement a whole family of languages targeting the ASM model
  - A new approach for concurrent ASM, which can be used to precisely describe distributed algorithms in ASM
  - optimizations can be plugged in, allowing enhancements of the generated code